# Smoothness Analysis for Probabilistic Programs with Application to Optimised Variational Inference



## Part 1 Smoothness Analysis for Probabilistic Programs with Application to Optimised Variational Inference





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Part 2

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"Smoothness" = {differentiability, Lipschitz continuity, continuity, ...}.

infinitely differentiable (in mathematics)

"Smoothness" = {differentiability, Lipschitz continuity, continuity, ...}.



Can apply many inference algorithms.

probabilistic model









"Smoothness" = {differentiability, Lipschitz continuity, continuity, ...}.

#### Goal: Find out **smoothness** properties<sup>?</sup> automatically and soundly.

 $= \mathbb{E}_{p(\boldsymbol{\epsilon})} \left[ \nabla_{\boldsymbol{\theta}} f(g(\boldsymbol{\epsilon}; \boldsymbol{\theta})) \right]$ 

of **programs** 

oathwise gradient estimator

 $\frac{di}{dt} = -\frac{1}{\partial q_i}$ Hamiltonian Monte Carlo expressed by programs

neural network

x

probabilistic model

Can provide provable robustness.

Can give guaranteed generalization bounds.

Smoothness = differentiability. Programs = deterministic, imperative programs.

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$$P \triangleq (y:=exp(x); if (x>0) \{z:=1\} else \{z:=-1\})$$
  
real-valued

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 $P \triangleq (y:=exp(x); if (x>0) \{z:=1\} else \{z:=-1\})$ 

• 
$$\llbracket P \rrbracket$$
 :  $\mathbb{R}^3 \to \mathbb{R}^3$   
 $\begin{pmatrix} x \\ y \\ z \end{pmatrix} \mapsto \begin{pmatrix} x' \\ y' \\ z' \end{pmatrix} = \begin{pmatrix} x \\ \exp(x) \\ \operatorname{sgn}(x) \end{pmatrix}$ 

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•  $P \vdash y$  is differentiable in x

$$\iff$$
 ...

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• P  $\vdash$  y is differentiable in x

 $\begin{pmatrix} x \\ y_0 \\ z \end{pmatrix} \stackrel{\llbracket P \rrbracket}{\mapsto} \begin{pmatrix} x' \\ y' \\ z' \end{pmatrix}$ 

for all  $y_0, z_0 \in \mathbb{R}$ .

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Smoothness = differentiability. Programs = deterministic, imperative programs.

 $P \triangleq (y:=exp(x); if (x>0) \{z:=1\} else \{z:=-1\})$ 

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• P  $\vdash$  y is differentiable in x

$$\Rightarrow f: \mathbb{R} \to \mathbb{R}$$
$$\begin{pmatrix} x \\ y_0 \\ z_0 \end{pmatrix} \stackrel{\mathbb{P}}{\mapsto} \begin{pmatrix} x' \\ y' \\ z' \end{pmatrix}$$

for all  $y_0, z_0 \in \mathbb{R}$ .

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•  $P \vdash y$  is differentiable in x

 $\Leftrightarrow f: \mathbb{R} \to \mathbb{R} \quad \text{is differentiable} \quad \text{for all } y_0, z_0 \in \mathbb{R}. \\ \begin{pmatrix} x \\ y_0 \\ z_0 \end{pmatrix} \stackrel{\mathbb{P}}{\mapsto} \begin{pmatrix} x' \\ y' \\ z' \end{pmatrix}$ 

Smoothness = differentiability. Programs = deterministic, imperative programs.

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• **P**  $\vdash$  y is differentiable in x.

. . .

**P**  $\nvDash$  z is differentiable in x.

Smoothness = differentiability. Programs = deterministic, imperative programs.

 $P \triangleq (y:=exp(x); if (x>0) \{z:=1\} else \{z:=-1\})$ 

It is surprisingly subtle to find out such smoothness properties (1) automatically, (2) soundly, and (3) precisely enough.

Want to check differentiability of a program in a compositional way.

Challenging case: Given P, P<sup>3</sup> and  $U, V \subseteq Var$ , want to check

 $\forall u \in U, \forall v \in V. \quad \mathsf{P}; \mathsf{P}' \vdash u \text{ is differentiable in } v$ .

Based on the chain rule:

 $\exists T \subseteq Var.$  $\forall u \in U, \forall t \in T.$  $P' \vdash u$  is differentiable in t $\forall t \in T, \forall v \in V.$  $P \vdash t$  is differentiable in v

• Looks sound by chain rule.

Seq

Based on the chain rule:

• Previously considered: e.g., [CACM'12] for continuity.

 $\exists T \subseteq \text{Var.} \quad \forall u \in U, \forall t \in T. \qquad P^{\prime} \vdash u \text{ is differentiable in } t \\ \forall t \in T, \forall v \in V. \qquad P \vdash t \text{ is differentiable in } v$ 

• Looks sound by chain rule.

Based on the chain rule:

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 $\exists T \subseteq \text{Var.} \quad \forall u \in U, \forall t \in T. \quad P' \vdash u \text{ is differentiable in } t \\ \forall t \in T, \forall v \in V. \quad P \vdash t \text{ is differentiable in } v \quad \checkmark \\ \text{Seq}$ 

 $\forall u \in U, \forall v \in V.$  P; P'  $\vdash u$  is differentiable in v

Is this rule indeed sound?

 $P;P' \triangleq (y:=sgn(x); z:=x+y).$ 

$$\begin{pmatrix} x \\ y \\ z \end{pmatrix} \xrightarrow{\mathbb{P}} \begin{pmatrix} x \\ \mathsf{sgn}(x) \\ z \end{pmatrix}, \begin{pmatrix} x \\ y \\ z \end{pmatrix} \xrightarrow{\mathbb{P}} \begin{pmatrix} x \\ y \\ z \end{pmatrix}$$

P;P'  $\triangleq$  (y:=sgn(x); z:=x+y).  $U \triangleq \{z\}, T \triangleq \{x\}, V \triangleq \{x\}$ .

 $\bigvee \forall u \in U, \forall t \in T. \qquad P' \vdash u \text{ is differentiable in } t$  $\forall t \in T, \forall v \in V. \qquad P \vdash t \text{ is differentiable in } v$ Seq

$$\begin{pmatrix} x \\ y \\ z \end{pmatrix} \xrightarrow{} \mathbb{P}^{\mathbb{I}} \begin{pmatrix} x \\ sgn(x) \\ z \end{pmatrix}, \quad \begin{pmatrix} x \\ y \\ z \end{pmatrix} \xrightarrow{} \mathbb{P}^{\mathbb{I}} \begin{pmatrix} x \\ y \\ x + y \end{pmatrix}$$

 $\mathsf{P};\mathsf{P}' \triangleq (\mathsf{y}:=\mathsf{sgn}(\mathsf{x}) ; \mathsf{z}:=\mathsf{x}+\mathsf{y}). \quad U \triangleq \{z\}, T \triangleq \{x\}, V \triangleq \{x\}.$ 

✓  $\forall u \in U, \forall t \in T.$ P' ⊢ u is differentiable in t
✓  $\forall t \in T, \forall v \in V.$ P ⊢ t is differentiable in v
Seq



 $\mathsf{P};\mathsf{P}' \triangleq (\mathsf{y}:=\mathsf{sgn}(\mathsf{x}) ; \mathsf{z}:=\mathsf{x}+\mathsf{y}). \quad U \triangleq \{z\}, T \triangleq \{x\}, V \triangleq \{x\}.$ 

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Seq

★  $\forall u \in U, \forall v \in V$ . P; P'  $\vdash u$  is differentiable in v



This rule is unsound! But why?

 $\mathsf{P};\mathsf{P}' \triangleq (\mathsf{y}:=\mathsf{sgn}(\mathsf{x}) ; \mathsf{z}:=\mathsf{x}+\mathsf{y}). \quad U \triangleq \{z\}, T \triangleq \{x\}, V \triangleq \{x\}.$ 

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 $\forall u \in U, \forall v \in V. P; P' \vdash u$  is differentiable in v



Seq

 $\mathsf{P};\mathsf{P}' \triangleq (\mathsf{y}:=\mathsf{sgn}(\mathsf{x}) ; \mathsf{z}:=\mathsf{x}+\mathsf{y}). \quad U \triangleq \{z\}, T \triangleq \{x\}, V \triangleq \{x\}.$ 

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 $\mathsf{P};\mathsf{P}' \triangleq (\mathsf{y}:=\mathsf{sgn}(\mathsf{x}) ; \mathsf{z}:=\mathsf{x}+\mathsf{y}). \quad U \triangleq \{z\}, T \triangleq \{x\}, V \triangleq \{x\}.$ 



 $\mathsf{P};\mathsf{P}' \triangleq (\mathsf{y}:=\mathsf{sgn}(\mathsf{x}); \mathsf{z}:=\mathsf{x}+\mathsf{y}). \quad U \triangleq \{z\}, T \triangleq \{x\}, V \triangleq \{x\}.$ 

 $\forall u \in U, \forall t \in T. \qquad P' \vdash u \text{ is differentiable in } t \\ \forall t \in T, \forall v \in V. \qquad P \vdash t \text{ is differentiable in } v \\ \hline \end{bmatrix}$ 

 $\forall u \in U, \forall v \in V.$  P; P'  $\vdash u$  is differentiable in v

Lesson: Need to consider dependency between variables.

non-differentiable us

#### Issue 1: Possible Fix

 $P;P' \triangleq (y:=sgn(x) ; z:=x+y). \quad U \triangleq \{z\}, \forall T \triangleq \{x\}, \forall V \triangleq \{x\}.$   $\forall u \in U, \forall t \in T! \quad P' \vdash u \text{ is differentiable in } t$   $\forall t \in T, \forall v \in V. \quad P \vdash t \text{ is differentiable in } v$   $\forall u \in U, \forall v \in V. \quad P;P' \vdash u \text{ is differentiable in } v$ 

$$\begin{pmatrix} x \\ y \\ z \end{pmatrix} \xrightarrow{} \mathbb{P} \begin{pmatrix} x \\ \operatorname{sgn}(x) \\ z \end{pmatrix} \xrightarrow{} \mathbb{P} \begin{bmatrix} x \\ \operatorname{sgn}(x) \\ \mathbb{P} \end{bmatrix} \xrightarrow{} \begin{pmatrix} x \\ \operatorname{sgn}(x) \\ x + \operatorname{sgn}(x) \end{pmatrix}$$

#### Issue 1: Possible Fix

$$P;P' \triangleq (y:=sgn(x) ; z:=x+y). \quad U \triangleq \{z\}, T \triangleq \{x\}, V \triangleq \{x\}.$$

$$\forall u \in U, \forall t \in T. \quad P' \vdash u \text{ is differentiable in } t$$

$$\forall t \in T, \forall v \in V. \quad P \vdash t \text{ is differentiable in } v$$

$$\forall u \in U, \forall v \in V. \quad P;P' \vdash u \text{ is differentiable in } v$$
Seq'



#### Issue 1: Possible Fix

$$P;P' \triangleq (y:=sgn(x) ; z:=x+y). \quad U \triangleq \{z\}, \forall T \triangleq \{x\}, \forall V \triangleq \{x\}.$$

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$$\forall u \in U, \forall v \in V. \quad P;P' \vdash u \text{ is differentiable in } v$$

$$Seq'$$

$$(x)$$

$$(x)$$

$$(y)$$

$$(x)$$

$$($$

This rule now looks sound. Are we done?

#### **Issue 2: Precision**

 $P;P' \triangleq (y:=sgn(x) ; z:=x+y). \quad U \triangleq \{x\}, T \triangleq \{x\}, V \triangleq \{x\}.$   $\forall u \in U, \forall t \in T. \quad P' \vdash u \text{ is differentiable in } t$   $\forall t \in T, \forall v \in V. \quad P \vdash t \text{ is differentiable in } v$   $\forall u \in U, \forall v \in V. \quad P;P' \vdash u \text{ is differentiable in } v$ Seq'

$$\begin{pmatrix} x \\ y \\ z \end{pmatrix} \xrightarrow{} \mathbb{P}^{\mathbb{I}} \begin{pmatrix} x \\ \operatorname{sgn}(x) \\ z \end{pmatrix} \xrightarrow{} \mathbb{P}^{\mathbb{I}} \begin{pmatrix} x \\ \operatorname{sgn}(x) \\ \mathbb{P}^{\mathbb{I}} \end{bmatrix} \begin{pmatrix} x \\ \operatorname{sgn}(x) \\ x + \operatorname{sgn}(x) \end{pmatrix}$$

#### **Issue 2: Precision**

$$P;P' \triangleq (y:=sgn(x) ; z:=x+y). \quad U \triangleq \{x\}, T \triangleq \{x\}, V \triangleq \{x\}.$$

$$\forall u \in U, \forall t \in T. \quad P' \vdash u \text{ is differentiable in } t$$

$$\forall t \in T, \forall v \in V. \quad P \vdash t \text{ is differentiable in } v$$

$$\forall u \in U, \forall u \in V. \quad P \vdash t \text{ is differentiable in } v$$

$$\forall u \in U, \forall u \in V. \quad P \vdash v \text{ is differentiable in } v$$

 $\nabla u \in U, \nabla v \in V$ . P;  $P' \vdash u$  is differentiable in v



#### **Issue 2: Precision**


$\forall u \in U, \ \forall t \in \text{Var. P'} \vdash u \text{ is differentiable in } t$  $\forall t \in \text{Var}, \forall v \in V. P \vdash t \text{ is differentiable in } v$  $\forall u \in U, \forall v \in V. P; P' \vdash u \text{ is differentiable in } v$ 

Admit the imprecision for now. Is this rule indeed sound?

P;P'  $\triangleq$  (y:=x; z:=f(x,y)) for  $f(x,y) \triangleq \begin{cases} xy/(x^2 + y^2) & \text{if } (x,y) \neq (0,0) \\ 0 & \text{if } (x,y) = (0,0). \end{cases}$ 

 $\forall u \in U, \ \forall t \in \text{Var. P'} \vdash u \text{ is differentiable in } t \\ \forall t \in \text{Var}, \forall v \in V. P \vdash t \text{ is differentiable in } v \\ \hline \text{Seq'}$ 

$$\begin{pmatrix} x \\ y \\ z \end{pmatrix} \xrightarrow{\mathbb{P}} \begin{pmatrix} x \\ x \\ z \end{pmatrix} \xrightarrow{\mathbb{P}} \begin{pmatrix} x \\ x \\ z \end{pmatrix} \xrightarrow{\mathbb{P}} \begin{pmatrix} x \\ y \\ z \end{pmatrix} \xrightarrow{\mathbb{P}} \begin{bmatrix} x \\ y \\ f(x, y) \end{pmatrix}$$

P;P'  $\triangleq$  (y:=x; z:=f(x,y)) for  $f(x,y) \triangleq \begin{cases} xy/(x^2 + y^2) & \text{if } (x,y) \neq (0,0) \\ 0 & \text{if } (x,y) = (0,0). \end{cases}$  $U \triangleq \{z\}, V \triangleq \{x\}.$ 

 $\bigvee \forall u \in U, \forall t \in Var. P' \vdash u \text{ is differentiable in } t$  $\forall t \in Var, \forall v \in V. P \vdash t \text{ is differentiable in } v$ Seq'



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✓  $\forall u \in U, \forall t \in Var. P' \vdash u$  is differentiable in t✓  $\forall t \in Var, \forall v \in V. P \vdash t$  is differentiable in vSeq'



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✓  $\forall u \in U, \forall t \in Var. P' \vdash u$  is differentiable in t✓  $\forall t \in Var, \forall v \in V. P \vdash t$  is differentiable in vSeq'

 $\forall u \in U, \forall v \in V.$  P; P'  $\vdash u$  is differentiable in v



This rule is still unsound! But why?

P;P'  $\triangleq$  (y:=x; z:=f(x,y)) for  $f(x,y) \triangleq \begin{cases} xy/(x^2 + y^2) & \text{if } (x,y) \neq (0,0) \\ 0 & \text{if } (x,y) = (0,0). \end{cases}$  $U \triangleq \{z\}, V \triangleq \{x\}.$ 

 $\checkmark \forall u \in U, \forall t \in Var. P' \vdash u \text{ is differentiable in } t$  $\forall t \in Var, \forall v \in V. P \vdash t \text{ is differentiable in } v$ Seq'





 $(\dots / (\dots 2 + \dots 2) \text{ if } (r \cdot v) \neq (0,0)$  $P;P' \triangleq (y:$ = (0,0). g, h are partially differentiable  $\implies g \circ h$  does so.  $U \triangleq \{z\}, V \triangleq$ assumed implicitly,  $\checkmark$   $\forall u \in U, \forall t \in Var. P' \vdash u$  is differentiable in t but invalid. ✓  $\forall t \in Var, \forall v \in V.$  P  $\vdash t$  is differentiable in v Seq'  $\forall u \in U, \forall v \in V.$  P; P'  $\vdash u$  is differentiable in v  $\begin{pmatrix} x \\ y \end{pmatrix} \begin{pmatrix} \text{id, id} \end{pmatrix} \begin{pmatrix} x \\ x \end{pmatrix} f \begin{pmatrix} x \\ 1[x \neq 0] \cdot \frac{1}{2} \end{pmatrix}$ partially differentiable



Lesson: Need to identify & check assumptions on target smoothness.



It is subtle to do smoothness analysis, soundly and precisely.

→ Our approach for smoothness analysis

- $\phi \subseteq \{f : \mathbb{R}^n \to \mathbb{R}^m\}$ : any set of "smooth" functions.
  - E.g., {*f* : partially differentiable}.

- $\phi \subseteq \{f : \mathbb{R}^n \to \mathbb{R}^m\}$ : any set of "smooth" functions.
  - E.g.,  $\{f: partially differentiable\}, \{f: jointly differentiable\}, \{f: continuous\}, \cdots$ .

- $\phi \subseteq \{f : \mathbb{R}^n \to \mathbb{R}^m\}$ : any set of "smooth" functions.
  - E.g.,  $\{f: partially differentiable\}, \{f: jointly differentiable\}, \{f: continuous\}, \cdots$ .
- P  $\vdash \{x, y, z\}$  is  $\phi$ -smooth in  $\{y, z\}$ 
  - $\iff \cdots$

- $\phi \subseteq \{f : \mathbb{R}^n \to \mathbb{R}^m\}$ : any set of "smooth" functions.
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- $\phi \subseteq \{f : \mathbb{R}^n \to \mathbb{R}^m\}$ : any set of "smooth" functions.
  - E.g.,  $\{f: partially differentiable\}, \{f: jointly differentiable\}, \{f: continuous\}, \cdots$ .
- $P \vdash \{x, y, z\} \text{ is } \phi \text{-smooth in } \{y, z\}$   $\iff f : \mathbb{R}^2 \rightarrow \mathbb{R}^3 \in \phi \text{ for any } x_0 \in \mathbb{R}.$  $\begin{pmatrix} x_0 \\ y \\ z \end{pmatrix} \stackrel{\text{IPI}}{\mapsto} \begin{pmatrix} x' \\ y' \\ z' \end{pmatrix}$





• Invariants:  $P \vdash V$  is dependent at most on  $d_P(V)$ .  $P \vdash V$  is  $\phi$ -smooth in  $s_P(V)$ .



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- Rules:  $d_{\mathrm{P};\mathrm{P}'}(V) \triangleq d_{\mathrm{P}}(d_{\mathrm{P}'}(V)).$



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• Rules:  $d_{\mathbf{P};\mathbf{P}'}(V) \triangleq d_{\mathbf{P}}(d_{\mathbf{P}'}(V)).$  $s_{\mathbf{P};\mathbf{P}'}(V) \triangleq s_{\mathbf{P}}(d_{\mathbf{P}'}(V)) \cap d_{\mathbf{P}}(s_{\mathbf{P}'}(V)^{c})^{c}.$  $u \in u \in \Lambda \quad u \in$ 









#### <u>Theorem</u> Our $\phi$ -smoothness analysis is sound if $\phi$ satisfies five assumptions:

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- ...
- ...
- ...
- ...

<u>Theorem</u> Our  $\phi$ -smoothness analysis is sound if  $\phi$  satisfies five assumptions:

- Composition:  $\forall f : \mathbb{R}^n \to \mathbb{R}^m, g : \mathbb{R}^m \to \mathbb{R}^l. \quad f, g \in \phi \implies (g \circ f) \in \phi.$
- Pairing:  $\forall f : \mathbb{R}^n \to \mathbb{R}^m, g : \mathbb{R}^n \to \mathbb{R}^l. \quad f, g \in \phi \implies (f, g) \in \phi.$
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- • •

<u>Theorem</u> Our  $\phi$ -smoothness analysis is sound if  $\phi$  satisfies five assumptions:

<ul> <li>Composition:</li> </ul>	$\forall f \colon \mathbb{R}^n \to \mathbb{R}^m$ , $g \colon \mathbb{R}^m \to \mathbb{R}^l$ .	$f, g \in \phi \implies$	$(g \circ f) \in \phi$

• Pairing: $\forall f : \mathbb{R}^n \to \mathbb{R}^m, g : \mathbb{R}^n \to \mathbb{R}^l.$	$f, g \in \phi \implies (f, g) \in \phi$	Ξφ,
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•••	Target smoothness property	A3 (proj.)	A4 (pair.)	A5 (rest.)	A6 (comp.)	A7 (stri.)
•••	cont. $(C^0)$	0	0	0	0	0
	locally Lipschitz (= $\phi^{(l)}$ )	0	0	0	0	0
•••	uniformly cont.	0	0	0	0	0
	Lipschitz cont.	0	0	0	0	0
	jointly diff.	0	0	0	0	0
	continuously diff. ( $C^1$ )	0	0	0	0	0
	smooth $(C^{\infty})$	0	0	0	0	0
	real analytic ( $C^{\omega}$ )	0	0	0	0	0
	partially cont. (= $\phi^{(pc)}$ )	0	0	0	×	0
	partially diff. (= $\phi^{(pd)}$ )	0	0	0	×	0
	almost-everywhere cont.	0	0	×	×	0
	almost-everywhere diff.	0	0	×	×	0
	coordinatewise non-decreasing	0	0	0	0	0
	locally bounded	0	0	0	0	0
	bounded	×	0	0	0	0
	Borol moogurable	0	0	0	0	0

<u>Theorem</u> Our  $\phi$ -smoothness analysis is sound if  $\phi$  satisfies five assumptions:

<ul> <li>Composition:</li> </ul>	$\forall f \colon \mathbb{R}^n \to \mathbb{R}^m, g \colon \mathbb{R}^m \to \mathbb{R}^l.$	$f, g \in \phi \implies$	$(g \circ f) \in \phi$

• Pairing:	$orall f \colon \mathbb{R}^n  o \mathbb{R}^m$ , $g \colon \mathbb{R}^n  o \mathbb{R}$	$\mathbb{R}^l$ . $f, g \in \phi$	$\implies (f,g) \in \phi.$
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•••	Target smoothness property	A3 (proj.)	A4 (pair.)	A5 (rest.)	A6 (comp.)	A7 (stri.)
•	cont. $(C^0)$	0	0	0	0	0
	locally Lipschitz (= $\phi^{(l)}$ )	0	0	0	0	0
• •••	uniformly cont.	0	0	0	0	0
	Lipschitz cont.	0	0	0	0	0
	jointly diff.	0	0	0	0	0
	continuously diff. ( $C^1$ )	0	0	0	0	0
	smooth ( $C^{\infty}$ )	0	0	0	0	0
	real analytic ( $C^{\omega}$ )	0	0	0	0	0
	partially cont. (= $\phi^{(pc)}$ )	0	0	0	×	0
	partially diff. (= $\phi^{(pd)}$ )	0	0	0	×	0
	almost-everywhere cont.	0	0	×	×	0
	almost-everywhere diff.	0	0	×	×	0
	coordinatewise non-decreasing	0	0	0	0	0
	locally bounded	0	0	0	0	0
	bounded	×	0	0	0	0
	Borol moogurable	0	0	0	0	0

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<u>Theorem</u> Our  $\phi$ -smoothness analysis is sound if  $\phi$  satisfies five assumptions:

- Composition:  $\forall f : \mathbb{R}^n \to \mathbb{R}^m, g : \mathbb{R}^m \to \mathbb{R}^l. \quad f, g \in \phi \implies (g \circ f) \in \phi.$
- Pairing:  $\forall f : \mathbb{R}^n \to \mathbb{R}^m, g : \mathbb{R}^n \to \mathbb{R}^l. \quad f, g \in \phi \implies (f, g) \in \phi.$

• • • •	Target smoothness property	A3 (proj.)	A4 (pair.)	A5 (rest.)	A6 (comp.)	A7 (stri.)
• •••	cont. $(C^0)$	0	0	0	0	0
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• •••	uniformly cont.	0	0	0	0	0
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	jointly diff.	0	0	0	0	0
	continuously diff. ( $C^1$ )	0	0	0	0	0
	smooth ( $C^{\infty}$ )	0	0	0	0	0
	real analytic ( $C^{\omega}$ )	0	0	0	0	0
	partially cont. (= $\phi^{(pc)}$ )	0	0	0	×	0
	partially diff. $(= \phi^{(pd)})$	0	0	0	×	0
	almost-everywhere cont.	0	0	×	×	0
	almost-everywhere diff.	0	0	×	X	0
	coordinatewise non-decreasing	0	0	0	0	0
	locally bounded	0	0	0	0	0
	bounded	×	0	0	0	0
	Borol magazzabla	0	0	0	0	0

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# Part 1 Smoothness Analysis for Probabilistic Programs with Application to Optimised Variational Inference





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Part 2

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# Variational Inference

• Problem: Given probability density functions p(z, x) and  $q_{\theta}(z)$ ,

minimize 
$$\mathcal{L}(\theta) \triangleq \mathbb{E}_{q_{\theta}(z)} \left[ \log \frac{q_{\theta}(z)}{p(z,x)} \right]$$
 over  $\theta \in \mathbb{R}^{n}$ .

• Typical approach: Apply a gradient descent algorithm.

$$\theta_{t+1} \coloneqq \theta_t - \eta \cdot \nabla_{\theta} \mathcal{L}(\theta_t).$$

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$$\theta_{t+1} \coloneqq \theta_t - \eta \cdot \nabla_{\theta} \mathcal{L}(\theta_t).$$
  
difficult to compute, so only estimate

#### **Gradient Estimators**

• Basic estimator (called score estimator):

 $\nabla_{\theta} \mathcal{L}(\theta) \approx \cdots$ 



• "Optimized" estimator (called pathwise gradient estimator):

$$\nabla_{\theta} \mathcal{L}(\theta) \approx \cdots$$

for  $z \sim r(-)$ .

#### **Gradient Estimators**

• Basic estimator (called score estimator):

$$\nabla_{\theta} \mathcal{L}(\theta) \approx \log \frac{q_{\theta}(z)}{p(z,x)} \times \nabla_{\theta}(\log q_{\theta}(z)) \quad \text{for } z \sim q_{\theta}(-).$$

• "Optimized" estimator (called pathwise gradient estimator):

$$\nabla_{\theta} \mathcal{L}(\theta) \approx \overline{\nabla_{\theta}} \left( \log \frac{q_{\theta}(t_{\theta}(z))}{p(t_{\theta}(z), x)} \right)$$

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#### **Gradient Estimators**

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**Requirements:**  $q_{\theta}(z)$  should be differentiable in  $\theta$ , ....

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Requirements:  $q_{\theta}(z)$  and p(z, x) should be differentiable in  $\theta$  and  $z, \dots$ .

In practice, we apply optimized estimator selectively to some of  $z_1, \dots, z_m$ . To do so soundly, we need to know differentiable parts of  $q_{\theta}(z)$  and p(z, x).

## **Gradient Estimators**

• Basic estimator (called score estimator):

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Requirements:  $q_{\theta}(z)$  should be differentiable in  $\theta$ , ....

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$$\nabla_{\theta} \mathcal{L}(\theta) \approx \nabla_{\theta} \left( \log \frac{q_{\theta}(t_{\theta}(z))}{p(t_{\theta}(z), x)} \right) \qquad \text{for } z \sim r(-).$$
Requirements:  $q_{\theta}(z)$  and  $p(z, x)$  should be diferentiable programs
Part 1 is used here.
In practice, we apply optimized estimator selectively to some of  $z_1, \cdots, z_m$ .
To do so soundly, we need to know differentiable parts of  $q_{\theta}(z)$  and  $p(z, x)$ .

PL for variational inference		# of "optimized" random var's			
(with neural nets, ···)					
Example in Pyro		Our optimizer Pyro's default optimizer			
	LUC	# rv (Sound) # rv (Sound) # rv (Unsound)			
Splitting normal	16				
··· (7 examples omitted)	•••				
Deep exponential family	105				
Deep Markov model	112				
Hidden Markov model	137				
Single-cell annotation	147				
Attend-infer-repeat	174				
Conditional VAE	205				



PL for variational inference (with neural nets, …)		# of "opt	# of "optimized" random var's			
Evample in Dure	LoC	Our optimizer	Pyro's default optimizer			
Example in Pyro		# rv (Sound)	# rv (Sound)	# rv (Unsound)		
Splitting normal	16	1	1	1		
··· (7 examples omitted)	•••	•••	•••	•••		
Deep exponential family	105	6	6	0		
Deep Markov model	112	1	1	0		
Hidden Markov model	137	2	2	0		
Single-cell annotation	147	3	3	0		
Attend-infer-repeat	174	1	1	1		
Conditional VAE	205	1	1	0		

PL for variational inference (with neural nets, …)	2	# of "optimized" random var's			
Example in Pyro		Our optimizer	Pyro's defa	ult optimizer	
	LUC	# rv (Sound)	# rv (Sound)	# rv (Unsound)	
Splitting normal	16	1	1	1	
··· (7 examples omitted)	•••	•••	•••		
Deep exponential family	105	6	6	0 due to branching	
Deep Markov model	112	1	1	0	
Hidden Markov model	137	2	2	0 due te div by 0	
Single-cell annotation	147	3	3		
Attend-infer-repeat	174	1	1	1	
Conditional VAE	205	1	1	0	



# **High-Level Messages**

- It is subtle to do smoothness analysis properly (automatically, soundly, precisely enough). One reason: Make assumptions on target smoothness, which are easily violated.
- There are some PL research opportunities for ML (which are less explored). This work: Static analysis for automatic planning of inference algorithms.

#### Thanks for your attention!

## Our Approach: Soundness

<u>Theorem</u> Our  $\phi$ -smoothness analysis is sound if  $\phi$  satisfies five assumptions:

 $\begin{array}{lll} \text{Strictness:} & \forall f = \lambda x. \bot : \mathbb{R}^n \to \mathbb{R}^m. & f \in \phi. \\ \text{Projection:} & \forall f = \operatorname{proj}_{n \to m} : \mathbb{R}^n \to \mathbb{R}^m. & f \in \phi. \\ \text{Restriction:} & \forall f : \mathbb{R}^n \to \mathbb{R}^m, x \in \mathbb{R}^k (k \le n). & f \in \phi \implies f(x, -) \in \phi. \\ \text{Composition:} & \forall f : \mathbb{R}^n \to \mathbb{R}^m, g : \mathbb{R}^m \to \mathbb{R}^l. & f, g \in \phi \implies (g \circ f) \in \phi. \\ \text{Pairing:} & \forall f : \mathbb{R}^n \to \mathbb{R}^m, g : \mathbb{R}^n \to \mathbb{R}^l. & f, g \in \phi \implies (f, g) \in \phi. \end{array}$ 

77 · · · · · ·	10/ 1)	A 4 / · · ·	A = ( )	1 ( )	A = ( )
Target smoothness property	A3 (proj.)	A4 (pair.)	A5 (rest.)	A6 (comp.)	A7 (stri.)
cont. $(C^0)$	0	0	0	0	0
locally Lipschitz (= $\phi^{(l)}$ )	0	0	0	0	0
uniformly cont.	0	0	0	0	0
Lipschitz cont.	0	0	0	0	0
jointly diff.	0	0	0	0	0
continuously diff. ( $C^1$ )	0	0	0	0	0
smooth ( $C^{\infty}$ )	0	0	0	0	0
real analytic ( $C^{\omega}$ )	0	0	0	0	0
partially cont. (= $\phi^{(pc)}$ )	0	0	0	×	0
partially diff. (= $\phi^{(pd)}$ )	0	0	0	×	0
almost-everywhere cont.	0	0	×	×	0

# Our Approach: Soundness

Target smoothness property	A3 (proj.)	A4 (pair.)	A5 (rest.)	A6 (comp.)	A7 (stri.)
$\overline{\operatorname{cont.}(\mathcal{C}^0)}$	0	0	0	0	0
locally Lipschitz (= $\phi^{(l)}$ )	0	0	0	0	0
uniformly cont.	0	0	0	0	0
Lipschitz cont.	0	0	0	0	0
diff. (= $\phi^{(d)}$ )	0	0	0	0	0
continuously diff. $(C^1)$	0	0	0	0	0
smooth ( $C^{\infty}$ )	0	0	0	0	0
real analytic ( $C^{\omega}$ )	0	0	0	0	0
partially cont. (= $\phi^{(pc)}$ )	0	0	0	×	0
partially diff. (= $\phi^{(pd)}$ )	0	0	0	×	0
almost-everywhere cont.	0	0	×	×	0
almost-everywhere diff.	0	0	×	$\times$	0
coordinatewise non-decreasing	0	0	0	0	0
locally bounded	0	0	0	0	0
bounded	$\times$	0	0	0	0
Borel measurable	0	0	0	0	0
locally integrable	0	0	×	$\times$	0
integrable	×	0	×	×	0

#### Reparametrization trick is incorrect in presence of discontinuity #2277

